

Exercise 1. Read §1-5 of the paper S. Artemov. ‘Explicit provability and constructive semantics,’ Bulletin of Symbolic Logic, vol. 7, No.1, pp. 1-36, 2001. The paper is also available on the website <http://web.cs.gc.cuny.edu/~sartemov>.

Exercise 2. For the following modal formulas find their proofs in S4 and then build realizations of these proofs in LP:

- a) $\Box A \rightarrow \Box(B \rightarrow A)$;
- b) $\Box A \rightarrow \Box(B \vee \Box A)$;
- c) $\Box A \vee \Box B \rightarrow \Box(\Box A \vee B)$.

Hint. Derive in S4 first, then try to mimic the S4-derivation in LP. Use ‘+’ to reconcile different proof terms by the same formulas, of needed. One more example: find an explicit version of $\Box A \rightarrow \Box(B \rightarrow \Box A)$. First, we derive this formula in S4:

1. $\Box A \rightarrow (B \rightarrow \Box A)$, a propositional axiom;
2. $\Box(\Box A \rightarrow (B \rightarrow \Box A))$, from 1, by Necessitation;
3. $\Box(\Box A \rightarrow (B \rightarrow \Box A)) \rightarrow (\Box \Box A \rightarrow \Box(B \rightarrow \Box A))$, modal distributivity axiom;
4. $\Box \Box A \rightarrow \Box(B \rightarrow \Box A)$, by Modus Ponens, from 2,3;
5. $\Box A \rightarrow \Box \Box A$, positive introspection axiom;
6. $\Box A \rightarrow \Box(B \rightarrow \Box A)$, from 4,5, by propositional logic.

Then we emulate this derivation in LP using common sense and some ingenuity.

1. $x:A \rightarrow (B \rightarrow x:A)$, a propositional axiom;
2. $c:(x:A \rightarrow (B \rightarrow x:A))$, by Constant Specification rule R1;
3. $c:(x:A \rightarrow (B \rightarrow x:A)) \rightarrow (!x:x:A \rightarrow (c!x):(B \rightarrow x:A))$, Application Axiom;
4. $!x:x:A \rightarrow (c!x):(B \rightarrow x:A)$, by Modus Ponens, from 2,3;
5. $x:A \rightarrow !x:x:A$, proof checking axiom;
6. $x:A \rightarrow (c!x):(B \rightarrow x:A)$, from 4,5, by propositional logic.

Exercise 3. Prove the deduction theorem for LP: if $\Gamma, A \vdash B$ then $\Gamma \vdash A \rightarrow B$.

Exercise 4. Prove Godel’s theorem of 1938: if $\text{LP} \vdash F$ then $\text{LP} \vdash p.F$ for some proof polynomial p .

Exercise 5. Consider the forgetful projection F^o of a LP-formula F obtained from F by replacing $t:G$ to $\Box G$ for any subformula $!t:G$. Show that LP^o is a subset of theorems of S4.

Exercise 6. Do you think LP^o actually equals S4? Any proof ideas?

Exercise 7. Let F be a modal tautology $\Box p \rightarrow \Box p$.

- a) find its realization in LP (i.e. a formula G derivable in LP such that G^O is F) OTHER than $t:p \rightarrow t:p$.
- b) find a realization of F in LP different from (a).